# Computationally efficient PageRank algorithm exploiting graph sparsity

### Dmitry Kamzolov

Scientific advisor: Yury Maximov, Alexander Gasnikov Moscow Institute of Physics and Technology Department of Control and Applied Mathematics

20 May 2015

#### Purpose of research

Checking theoretical estimation algorithm based on the Nesterov ideas ranking web-pages on the sparse graphs.

### Application

It lets solve a large class ranking problems in logarithmic largest space time.

## Notation

- Given a directed graph with *n* vertices.
- The vertices sites.
- Oriented edge links.
- Matrix **P** is adjacency matrix for the graph.
- Vector PageRank x is vector quantity characterizing the importance of each site.
- Sparsity coefficient *s* is the maximum number of non-zero elements in each column and each row of the matrix **P**.



The problem of finding the vector PageRank = The problem of searching left eigenvector x, that

$$\mathbf{x}^T = \mathbf{x}^T \mathbf{P}$$
, where  $\sum_{k=1}^n x_k = 1$ .

The problem of searching left eigenvector = The optimization problem

$$f(\mathbf{x}) = \frac{1}{2} \|\mathbf{A}\mathbf{x}\|_2^2 + \frac{1}{2} \sum_{k=1}^n (-x_k)_+^2 \longrightarrow \min_{\langle \mathbf{x}, \mathbf{e} \rangle = 1},$$

where matrix  $\mathbf{A} = \mathbf{P}^T - \mathbf{I}$ ,  $\mathbf{I}$  - the identity matrix,  $\mathbf{e} = (1, \dots, 1)$ ,

- Matrix A is sparse with Lipschitz constant L.
- Componentwise descent.

$$x_{k+1} = x_k + \operatorname*{argmin}_{h:\langle h, e \rangle = 0} \left\{ f(x_k) + \langle \nabla f(x_k), h \rangle + \frac{L}{2} \|h\|_1^2 \right\},$$

- The number of iterations  $O(1/\epsilon^2)$
- The complexity of a single iteration  $O(s^2 \ln n)$
- The complexity of the algorithm

$$O\left(\frac{s^2\ln n}{\varepsilon^2}\right)$$

- The goal of computational experiment is to check theoretical estimates of the complexity.
- Input: the dimension n, the sparsity coefficient s, the generated sparse matrix **A** Erdos-Renyi type with parameters n, s, the initial value of the function  $f_0$  and the initial value of the gradient of  $g_0$ , precounted  $\mathbf{A}^T \mathbf{A}$ .
- Output: Page Rank x, time.

- Checking dependence  $O(1/\epsilon^2)$ . Timelines execution of the accuracy of the solution (1-2). Charts have a linear dependence.  $\implies$  The complexity  $O(1/\epsilon^2)$ . Theoretical estimates are confirmed.
- The checking of dependence.  $O(s^2 \ln n)$ . Unable to use the sparsity of the matrix due to the nature of MatLab and memory arrays. Theoretical estimates have not been confirmed.
- Getting Dependence O(n). Timeline of execution of the dimension (3). Charts have a linear dependence.  $\implies$  complexity O(n).
- The final complexity:  $O(n/\epsilon^2)$



pic.1 The dependence of the runtime by the number of iterations of the algorithm for n = 12500 sites



Рис.2 The dependence of the runtime by the number of iterations of the algorithm for n = 15000 sites



Рис.3 The dependence of the the runtime from the dimension n where 21000 iteration

999

æ

-

< / □ > <

It received the complexity:

- $O\left(\frac{n}{\epsilon^2}\right)$ , in practice without sparsity,
- $O\left(\frac{s^2 \ln n}{\varepsilon^2}\right)$  in theory with sparsity.
- The table shows the rate of convergence of the most advanced fast algorithms.
- Our algorithm has the lowest the complexity.

Method	Condition	Complexity
Nazin-Polyak 2008	no	$O\left(\frac{n\ln(\frac{n}{\sigma})}{\epsilon^2}\right)$
Nesterov 2012	S	$O\left(\frac{sn\ln n}{\epsilon^2}\right)$
Juditsky et al 2009	no	$O\left(\frac{n\ln(\frac{n}{\sigma})}{\epsilon^2}\right)$
Grigoriadis-Hachiyan 2009	5	$O\left(\frac{s\ln n\ln(\frac{n}{\sigma})}{\epsilon^2}\right)$
Polyak-Tremba 2012	S	$\frac{2sn}{\epsilon}$
Gasnikov-Dvurechensky 2015	S	$O\left(\frac{s^2 \ln n}{\epsilon^2}\right)$

Estimates of the rate of convergence of algorithms PageRank, where  ${\cal S}$  sparsity-condition.

In our work:

- The research method of ranking web pages with sparse graphs.
- It is shown that the theoretical estimate of the number of steps corresponds to the experimental data.
- It is shown that the theoretical estimate of the complexity of the algorithm step does not correspond to the experimental data, due to the nature of programmatic implementation.
- Through the use of 1-norm achieved O(n) arithmetic operations on a step of the algorithm.
- The result is a new, and even with the deterioration estimates of complexity algorithm is one of the fastest algorithms.