

On the relation between Kantorovich's and Shannon's optimization problems

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Higher School of Economics, Nizhny Novgorod, Russia

Optimal transportation problems (OTPs)

Information and entropy

Optimal channel problem (OCP)

Dual formulation of OTP

Geometry of information divergence and optimization

Relation between $K_c[p, q]$ and $D_{KL}[p, q]$

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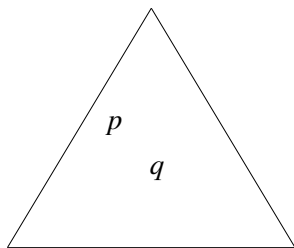
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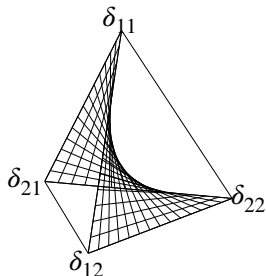
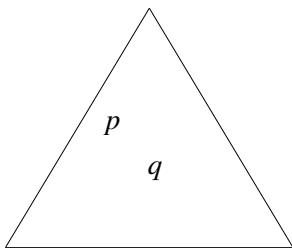
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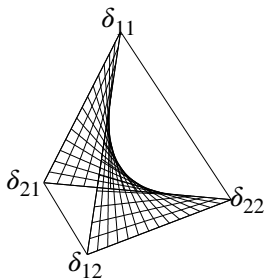


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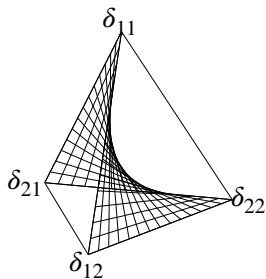
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$$T : \mathcal{P}(X) \rightarrow \mathcal{P}(Y):$$

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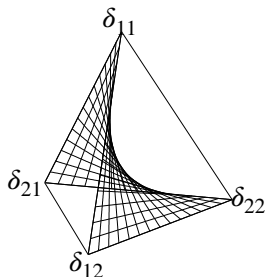
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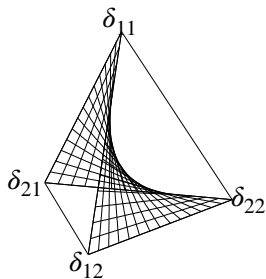
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- $p(y | x)$ — Markov transition kernel
- T is determined by $w \in \mathcal{P}(X \otimes Y)$:

$$w = p(y | x) \otimes q$$



Monge OTP

Optimal Transportation Problem (Monge, 1781)

$$K_c[p, q] := \inf \left\{ \int_X c(x, f(x)) dq : f : p = q \circ f^{-1} \right\}$$

where $p = q \circ f^{-1}$ is push-forward under measurable mapping $f : X \rightarrow Y$:

$$p(E) = q \circ f^{-1}(E) = q\{x : f(x) \in E\}$$

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Optimal Transport

- $p(E | x)$ has the form:

$$\delta_{f(x)}(E) = \begin{cases} 1 & \text{if } f(x) \in E \\ 0 & \text{otherwise} \end{cases}$$

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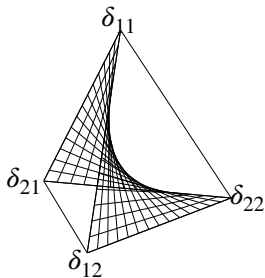
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- $w_f \in \partial \mathcal{P}(X \otimes Y)$:

$$w_f(X, Y \setminus f(X)) = 0$$



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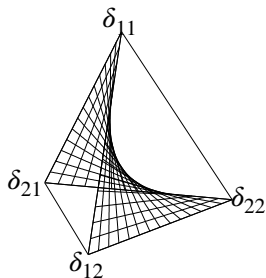
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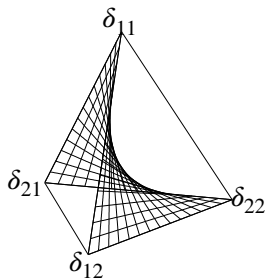
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For $w \in \Gamma[q, p] \subset \mathcal{P}(X \otimes Y)$:

$$\begin{aligned} I_w\{x, y\} &:= D_{KL}[w, q \otimes p] \\ &= H[q] - H[q(x | y)] \\ &= H[p] - H[p(y | x)] \end{aligned}$$



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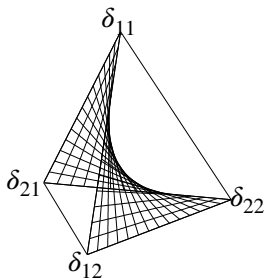
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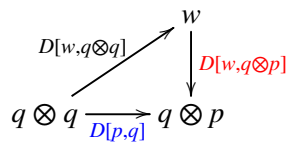
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Entropy $H[p] = - \int \ln p dp$

$$H[p] := I_w\{y, y\} = \sup_{w: \pi_Y w = p} I_w\{x, y\}$$



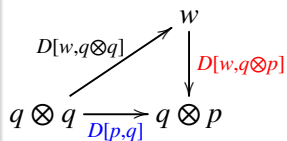


Theorem (Shannon-Pythagorean)

- $w \in \mathcal{P}(X \otimes Y)$, $\pi_X w = q$, $\pi_Y w = p$

$$D_{KL}[w, q \otimes q] = D_{KL}[w, q \otimes p] + D_{KL}[p, q]$$

(Belavkin, 2013a)

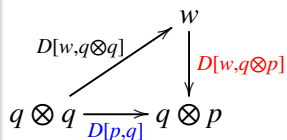


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Proof.

$$D[w, q \otimes q] = \underbrace{D[w, q \otimes p]}_{I_w\{x,y\}} + \underbrace{D[q \otimes p, q \otimes q]}_{D[p,q]} - \underbrace{\langle \ln q \otimes p - \ln q \otimes q, q \otimes p - w \rangle}_0$$

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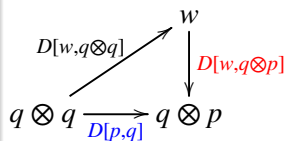
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Cross-Information (Belavkin, 2013a)

$$D_{KL}[w, q \otimes q] = \underbrace{-\langle \ln q, p \rangle}_{\text{Cross-entropy}} - \underbrace{\left(H[p] - D_{KL}[w, q \otimes p] \right)}_{H[p(y|x)]}$$

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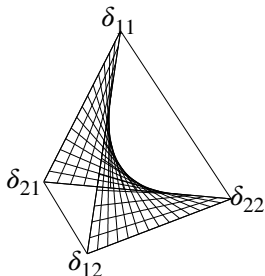
Optimal Channel Problem (Shannon, 1948)

$$R_c[q](\lambda) := \inf \left\{ \int_{X \times Y} c(x, y) d\omega : \pi_X \omega = q, I_\omega\{x, y\} \leq \lambda \right\}$$

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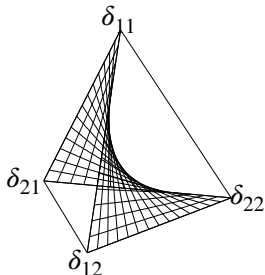
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Exponential family solutions

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$$w = e^{-\beta c - \ln Z} q \otimes p, \quad \beta^{-1} = -dR_c[q](\lambda)/d\lambda$$



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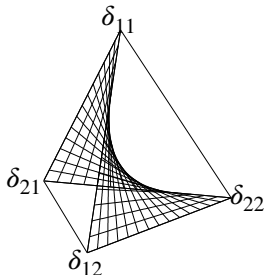
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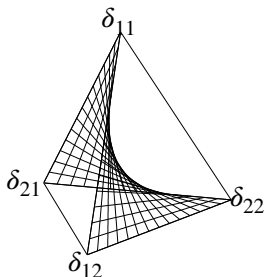
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Value of Information (Stratonovich, 1965)

$$V(\lambda) := R_c[q](0) - R_c[q](\lambda) = \sup \{ \mathbb{E}_w \{u\} : I_w\{x, y\} \leq \lambda \}$$

Relation to Kantorovich OTP

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$$0 \leq I_w\{x, y\} \leq \min\{H[q], H[p]\}$$

- $K_c[q, p]$ has implicit constraint $I_w\{x, y\} \leq \lambda = \min\{H[q], H[p]\}$ and

$$R_c[q](\lambda) \leq K_c[q, p](\lambda)$$

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Inverse of the OCP Value

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- These inverse values represent the smallest amount of Shannon's information required to achieve expected cost $\int c dw = v$.
- If $v = K_c[q, p](\lambda)$, then

$$R_c^{-1}[q](v) \leq K_c^{-1}[q, p](v)$$

Common Solution

Theorem

Let w_{OCP} and $w_{OTP} \in \mathcal{P}(X \times Y)$ be optimal solutions to OCP and OTP problems with the same constraint $I(x, y) \leq \lambda$. Then $R_c[q](\lambda) = K_c[p, q](\lambda)$ if and only if $w_{OCP} = w_{OTP} \in \Gamma[p, q]$.

Common Solution

Theorem

Let w_{OCP} and $w_{OTP} \in \mathcal{P}(X \times Y)$ be optimal solutions to OCP and OTP problems with the same constraint $I(x, y) \leq \lambda$. Then $R_c[q](\lambda) = K_c[p, q](\lambda)$ if and only if $w_{OCP} = w_{OTP} \in \Gamma[p, q]$.

Proof.

- $w_{OCP} \in \partial D^*[-\beta c, q \otimes p]$ of subdifferential of D^* at $u = -\beta c$:

$$D^*[u, q \otimes p] = \ln \int_{X \times Y} e^{u(x,y)} dq(x) dp(y)$$

- $w_{OTP} \in \partial D^*[-\beta c, q \otimes p]$ implies

$$(1-t)w_{OCP} + tw_{OTP} \in \partial D^*[-\beta c, q \otimes p], \quad t \in [0, 1]$$

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- and the KL-divergence $D[w, q \otimes p]$, the dual of $D^*[u, q \otimes p]$, not strictly convex.

Optimal transportation problems (OTPs)

Information and entropy

Optimal channel problem (OCP)

Dual formulation of OTP

Geometry of information divergence and optimization

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- Consider $f : X \rightarrow \mathbb{R}$ and $g : Y \rightarrow \mathbb{R}$ such that

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- Can $J_c[p, q]$ be related to $D_{KL}[p, q]$?

KL-divergence



$$\begin{aligned}
 D[p, q] &= D[p, r] + D[r, q] - \int_X \ln \frac{dq(x)}{dr(x)} [dp(x) - dr(x)] \\
 &= D[p, r] - D[q, r] - \int_X \ln \frac{dq(x)}{dr(x)} [dp(x) - dq(x)]
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- Let us denote

$$J_{c,\varepsilon}[p, q] := \frac{1}{\varepsilon} [\beta \mathbb{E}_p\{f\} - \alpha \mathbb{E}_q\{g\}]$$

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Common solution for Dual OTP

Theorem

Let (f, g) be the solution to the dual OTP. If there exists a reference measure $r \in \mathcal{P}(X)$ such that $f = \nabla D[p, r]$ and $g = \nabla D[q, r]$, then

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Proof.

- $f = \nabla D[p, r]$ and $g = \nabla D[q, r]$ imply that $\alpha = \beta = 1$, and

$$p = \exp(f - \kappa[f]) r, \quad q = \exp(g - \kappa[g]) r$$

- The result follows, because

$$\mathbb{E}_p\{f\} - \mathbb{E}_q\{g\} = J_c[p, q] = K_c[p, q]$$

Optimal transportation problems (OTPs)

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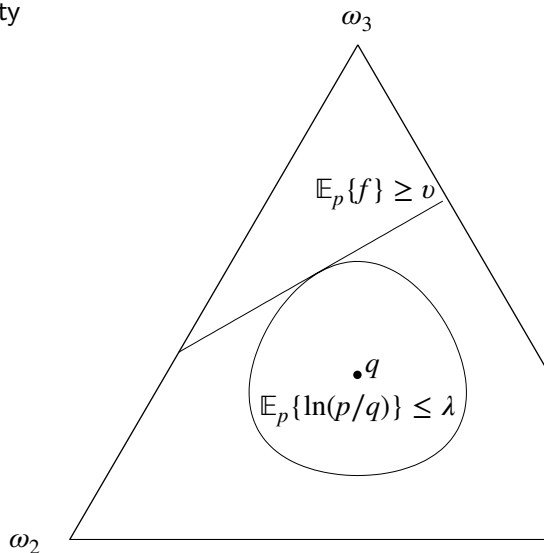
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Problems on Conditional Extremum

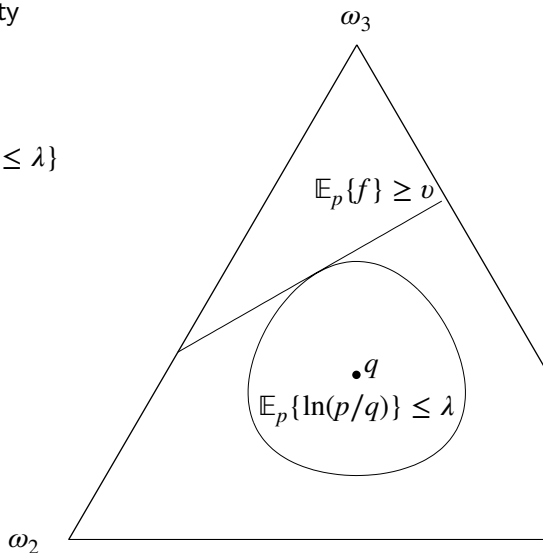
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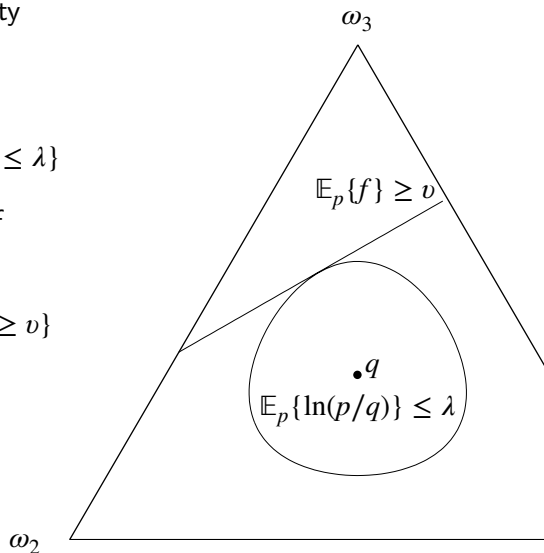
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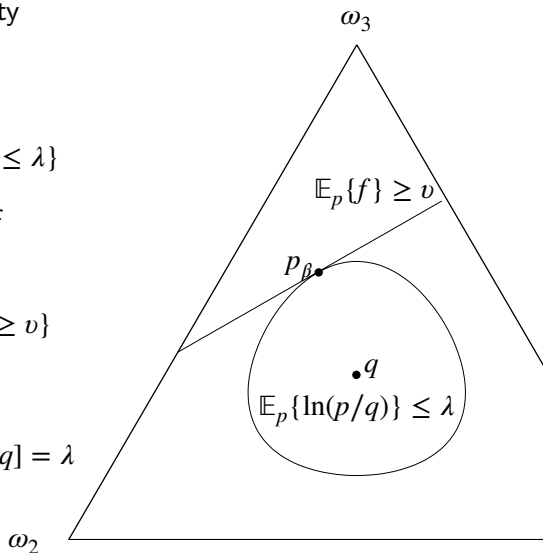
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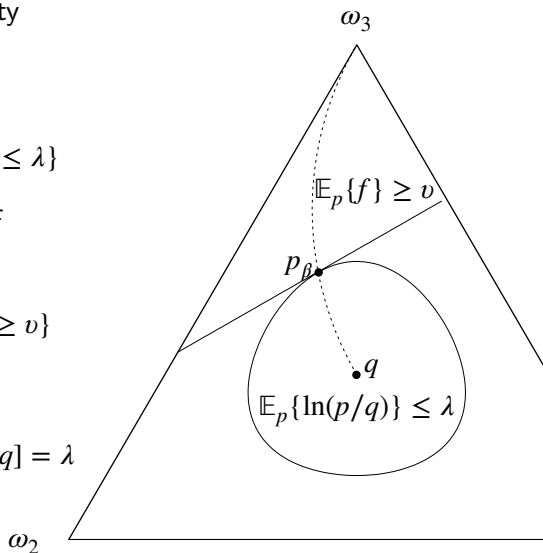
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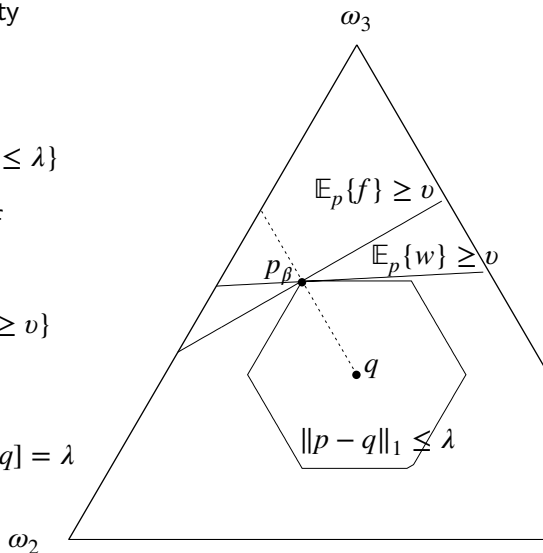
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General Solution

- Lagrange function for $v_u(\lambda) := \sup\{\langle u, p \rangle : F[p, q] \leq \lambda\}$
 $(\lambda_u(v) := \inf\{F[p, q] : \langle u, p \rangle \geq v\})$:

$$L(p, \beta^{-1}) = \langle u, p \rangle + \beta^{-1}(\lambda - F[p, q])$$

$$\left(L(p, \beta) = F[p, q] + \beta(v - \langle u, p \rangle) \right)$$

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- Optimal solutions are subgradients of $F^*[u, q] = \sup\{\langle u, p \rangle - F[p, q]\}$:

$$p(\beta) \in \partial F^*[\beta u], \quad F[p, q] = \lambda \quad \left(\langle u, p(\beta) \rangle = v \right)$$

Example: Exponential Solution

- For $F[p, q] = D_{KL}[p, q]$:

$$L(p, \beta^{-1}) = \langle u, p \rangle + \beta^{-1}(\lambda - \langle \ln(p/q), p \rangle + \langle 1, p - q \rangle)$$

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- Optimal solutions are gradients of $D_{KL}^*[u, q] = \ln \langle e^u, q \rangle$:

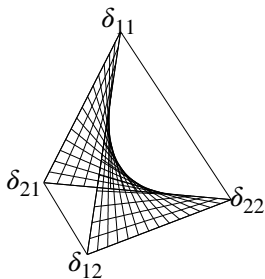
$$p(\beta) = e^{\beta u - \ln Z[\beta u]} q, \quad D_{KL}[p(\beta), q] = \lambda$$

Solution to Shannon's OCP

- The solution for

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$$w(\beta) = e^{\beta u - \ln Z[\beta u]} q \otimes p$$



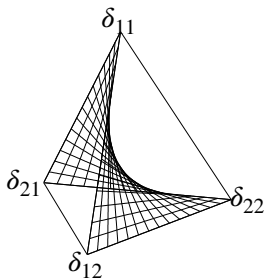
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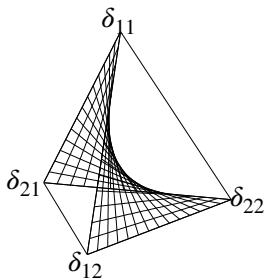
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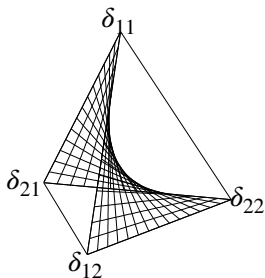
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- The dual is strictly convex:

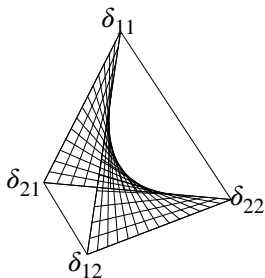
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Solution to Kantorovich's OTP

- $\Gamma[q, p]$ is convex:

$$\pi_X[(1-t)w_1 + tw_2] = (1-t)q + tq = q$$



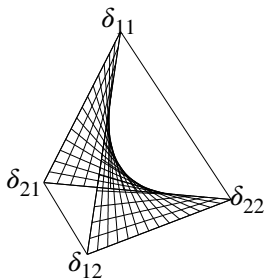
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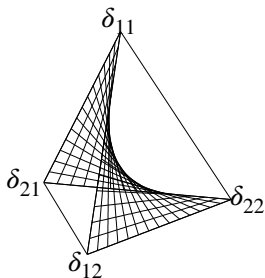
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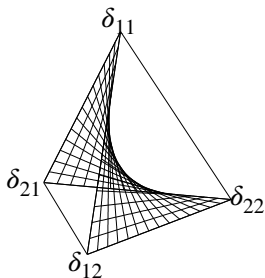
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Monge-Ampere equation

$$q = p \circ \nabla \varphi |\nabla^2 \varphi|$$

where $\varphi : X \rightarrow \mathbb{R} \cup \{\infty\}$ is convex, and $\nabla \varphi : X \rightarrow Y$ is such that $p = q \circ (\nabla \varphi)^{-1}$ (McCann, 1995; Villani, 2009).

Strict Inequalities

Theorem (Belavkin, 2013b)

- Let $\{w(\beta)\}_u$ be a family of $w(\beta) \in \mathcal{P}(X \otimes Y)$

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- $w(\beta) \in \partial \mathcal{P}(X \otimes Y)$ iff $\lambda \geq \sup F$ (i.e. $\beta \rightarrow \infty$).
- For any $\sigma \in \partial \mathcal{P}(X \otimes Y)$ with $F[\sigma] = F[w(\beta)] = \lambda$

$$\mathbb{E}_\sigma\{u\} < \mathbb{E}_{w(\beta)}\{u\}$$

Strict Inequalities

Theorem (Belavkin, 2013b)

- Let $\{w(\beta)\}_u$ be a family of $w(\beta) \in \mathcal{P}(X \otimes Y)$

maximizing $\mathbb{E}_w\{u\}$ on sets $\{w : F[w] \leq \lambda\}$, $\forall \lambda = F[w]$

- $F : \mathcal{P}(X \otimes Y) \rightarrow \mathbb{R} \cup \{\infty\}$ closed convex and minimized at

$$q \otimes p \in \partial F^*(0) \subset \text{Int}(\mathcal{P}(X \otimes Y))$$

- If F^* is strictly convex, then

- $w(\beta) \in \partial \mathcal{P}(X \otimes Y)$ iff $\lambda \geq \sup F$ (i.e. $\beta \rightarrow \infty$).
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$$\mathbb{E}_\sigma\{u\} < \mathbb{E}_{w(\beta)}\{u\}$$

- For any $\sigma \in \partial \mathcal{P}(X \otimes Y)$ with $\mathbb{E}_\sigma\{u\} = \mathbb{E}_{w(\beta)}\{u\} = v$

$$F[\sigma] > F[w(\beta)]$$

Strict Bounds for Monge OTP

Corollary

- *Let $w_f \in \Gamma[q, p]$ be a solution to Monge OTP $K_c[p, q]$.*

Strict Bounds for Monge OTP

Corollary

- Let $w_f \in \Gamma[q, p]$ be a solution to Monge OTP $K_c[p, q]$.
- Let $w(\beta)$ is a solution to Shannon's OCP $R_c[q](\lambda)$.

Strict Bounds for Monge OTP

Corollary

- Let $w_f \in \Gamma[q, p]$ be a solution to Monge OTP $K_c[p, q]$.
- Let $w(\beta)$ is a solution to Shannon's OCP $R_c[q](\lambda)$.
- If w_f and $w(\beta)$ have equal $I_{w_f}\{x, y\} = I_{w(\beta)}\{x, y\} = \lambda < \sup I_w\{x, y\}$, then

$$K_c[p, q](\lambda) > R_c[q](\lambda) > 0$$

Strict Bounds for Monge OTP

Corollary

- Let $w_f \in \Gamma[q, p]$ be a solution to Monge OTP $K_c[p, q]$.
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- If w_f and $w(\beta)$ achieve equal values $K_c[p, q] = R_c[q](\lambda) = v > 0$, then

$$K_c^{-1}[p, q](v) > R_c^{-1}[q](v)$$

Optimal transportation problems (OTPs)

Information and entropy

Optimal channel problem (OCP)

Dual formulation of OTP

Geometry of information divergence and optimization

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